

Distributed Systems

4. Naming and Binding

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- My 15" MacBook Pro
- The rightmost computer on my desk
- Paul's aluminum laptop, but not the big or the small one
- hedwig
- hedwig.pk.org
- 192.168.60.148
- 00:14:51:ec:f2:5b

Naming things

- User names
 - Login, email
- Machine names
 - rlogin, email, web
- Files
- Devices
- Variables in programs
- Network services

Naming Service

Allows you to look up **names**

- Often returns an *address* as a response

Might be implemented as

- Search through file
- Client-server program
- Database query
- ...

What's a name?

Name: identifies what you want

Address: identifies where it is

Route: identifies how to get there

Binding: associates a name with an address

- "choose a lower-level-implementation for a higher-level semantic construct"

RFC 1498: Inter-network Naming, addresses, routing

Names

We may need names for:

- **Services:** e.g., time of day
- **Nodes:** computer that can run services
- **Paths:** route
- **Objects within service:** e.g. files on a file server

Naming convention can take any format

- Ideally one that will suit application and user
- E.g., human readable names for humans, binary identifiers for machines

Uniqueness of names

- Easy on a small scale
- Problematic on a large scale
 - Globally unique names are hard to keep unique
- Hierarchy allows uniqueness to be maintained
 - Globally unique components
 - E.g.:
 - Ethernet MAC address: 3 bytes: organization, 3 bytes: controller
 - IP address: network & host
 - **compound name**: set of atomic names connected with a name separator
 - Domain name (www.rutgers.edu), URLs

Terms: Naming convention

Naming system determines syntax for a name

- Unix file names:
 - Parse components from left to right separated by /
 - /home/paul/src/gps/gui.c
- Internet domain names:
 - Ordered right to left and delimited by .
 - www.cs.rutgers.edu
- LDAP names
 - Attribute/value pairs ordered right to left, delimited by ,
 - cn=Paul Krzyzanowski, o=Rutgers, c=US

Terms: Context

- A particular set of name → object bindings
- Each context has an associated naming convention
- A name is *always* interpreted relative to some context
 - E.g., directory /usr in a UNIX file system

Terms: Naming System

Connected set of contexts of the same type (same naming convention) along with a common set of operations

For example:

- System that implements DNS
- System that implements LDAP

Terms: Name space

Set of names in the naming system

For example,

- Names of all files and directories in a UNIX file system
- All domain names on the Internet

Terms: Resolution

Name lookup

- Return the underlying representation of the name

For example,

- www.rutgers.edu → 128.6.4.5

Directory Service

Extension of naming service:

- Associates names with objects
- Allows objects to have attributes
- Can search based on attributes

For example,

- LDAP (Lightweight Directory Access Protocol)
- Directory can be an object store:
 - Look up printer object and send data stream to it

Name resolution

To send data to a service:

1. Find a node on which the service resides (service name resolution)
2. Find an address (or network attachment point) for that node (node name location)
3. Find a path from this location to the service (routing service)

Name resolution

E.g., access "paul's service":

File lookup: (service name to service address)

"paul's service" → cs.rutgers.edu:1234

DNS lookup: (domain name to IP address)

cs.rutgers.edu → 128.6.4.2

ARP resolution: (IP name to MAC address)

128.6.4.2 → 08:00:20:90:9c:23

IP routing: (dynamic for IP)

route: remus → lcsr-gw → aramis

Binding

The association of a resolution

Static binding

- Hard-coded

Early binding

- Look up binding before use
- Cache previously used binding

Late binding

- Look up just before use

IP Domain Names

Human readable names

e.g. aramis.rutgers.edu

Hierarchical naming scheme

- No relation to IP address or network class

Example: DNS

Internet Domain Name Service

- Maps machine names (www.rutgers.edu) to IP addresses (128.6.4.5)

In the past:

- Search `/etc/hosts` for machine name
- File periodically downloaded from Network Information Center (NIC) at the Stanford Research Institute (SRI)

Internet Domain Name Space

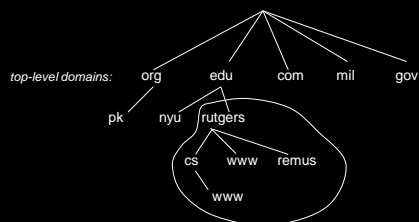
Tree structure

- Each node has resource information associated with it
- **owner**: domain name whose resource record is found
- **type of resource**:
 - Host address (A)
 - Alias name (C)
 - Name server for domain (NS)
 - Mail server (MX)
- **TTL** (time to live) → for **caching**
- **Relevant data** (e.g., address)

Domain Name Server

Essential task

Answer queries about data in its **zone**
(group of machines under a root – e.g. rutgers)



Sample Query

- Rutgers registers **rutgers.edu** with a domain registry
 - educause.net for .edu domain
 - See internic.net for ICANN-accredited list of registrars for top-level domains
- Top-level domain names and their associated name server info loaded to **root name servers**
 - 13 computers: replicated information
 - Contain addresses for all registries of top-level domains (.com, .edu, .org, ...)

Sample Query

Submit query to a local **DNS resolver**:

1. query(cs.rutgers.edu) → root name server
root name servers identify authoritative servers for top-level domains
send query to A.ROOT_SERVERS.NET: 198.41.0.4
2. referral to edu name server
returns list of DNS servers for .edu:
L3.NSTLD.COM: 192.41.162.32
3. query(cs.rutgers.edu) → edu name server
send query to 192.41.162.32
4. referral to rutgers.edu name servers:
 - DNS1.rutgers.edu 165.230.144.131
 - DNS2.rutgers.edu 128.6.21.9
 - DNS3.rutgers.edu 198.151.130.254
5. query(cs.rutgers.edu) → rutgers name server
send query to 165.230.144.131
6. rutgers name server returns
A: 128.6.4.2 address
MX: dragon.rutgers.edu domain name for email

Iterative vs. Recursive name resolution

- Referrals force iterative name resolution
 - Send query to root name server
 - Send query to edu name server
 - Send query to rutgers name server
 - *Advantage: stateless*
- Recursive:
 - Send query to root name server
 - Root name server sends query to edu name server
 - Edu name server sends query to rutgers name server
 - *Advantage: increased caching opportunities, reduced communication*

DNS

BIND

- Implementation of DNS provided by the Internet Software Consortium (www.isc.org)

Programs to perform queries:

- dnsquery, nslookup, dig, host

Naming: files

File system maps file pathname

`/home/paul/src/map.c`
↓ *namei* in kernel
`major=3, minor=6, inode=6160`

The End